# **ECE 8201: Low-dimensional Signal Models for High-dimensional Data Analysis**

Lecture 7: Matrix completion

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- "Guaranteed Minimum-Rank Solutions of Linear Matrix Equations via Nuclear Norm Minimization", Recht, Fazel, and Parrilo, 2007.
- "The power of convex relaxation: Near-optimal matrix completion", E. J. Candès and T. Tao, 2007.

## Outline

Matrix completion

- Motivation
- Theoretical aspects:
  - nuclear norm
  - low-rank matrix sensing
  - low-rank matrix completion
- efficient algorithm

The Netflix problem, or collaborative filtering

• How to estimate the missing ratings?



• About a million users, and 25,000 movies, with sparsely sampled ratings

• Matrix completion problem: consider  $M \in \mathbb{R}^{n_1 \times n_2}$  to represent the Netflix data set, we may model it through factorization:



• The rank r of M is much smaller than its dimension  $r \ll \min\{n_1, n_2\}$ .

- Given n points  $\{m{x}_j\}_{j=1}^n \in \mathbb{R}^3$
- Observe partial information about distances:

$$M_{i,j} = \| \boldsymbol{x}_i - \boldsymbol{x}_j \|_2^2$$

e.g. in wireless sensor network, each sensor can measure the distance to its neighbors, would like to globally locate all sensors.



• Write the matrix

$$oldsymbol{X} = egin{bmatrix} oldsymbol{x}_1^T \ oldsymbol{x}_2^T \ dots \ oldsymbol{x}_n^T \end{bmatrix} \in \mathbb{R}^{n imes 3}$$

then

$$M_{i,j} = \boldsymbol{x}_i^T \boldsymbol{x}_i + \boldsymbol{x}_j^T \boldsymbol{x}_j - 2\boldsymbol{x}_i^T \boldsymbol{x}_j$$

• Matrix completion problem: Let  $Y = XX^T$ . The distance matrix  $M \in \mathbb{R}^{n \times n}$  between points can be written as

$$M = \operatorname{diag}(Y)e^{T} + e\operatorname{diag}(Y)^{T} - 2Y$$

• The rank r of M is much smaller than its dimension  $r \ll n$ .

Structure from motion: reconstruct scene geometry and camera motion from multiple images.



In the pipeline of performing SFM, assume we've found a set of good feature points with their corresponding 2D locations in the images.



Tomasi and Kanade's factorization: Given n points  $x_{i,j}^T \in \mathbb{R}^2$  corresponding to the location of the *i*th point in the *j*th frame, define the matrix

$$oldsymbol{M} = egin{bmatrix} oldsymbol{x}_{1,1} & \cdots & oldsymbol{x}_{1,m} \ dots & \ddots & dots \ oldsymbol{x}_{n,1} & \cdots & oldsymbol{x}_{n,m} \end{bmatrix} \in \mathbb{R}^{n imes 2m}$$

In the absence of noise, this matrix admits a low-rank factorization:

$$\boldsymbol{M} = \underbrace{\begin{bmatrix} \boldsymbol{s}_1^T \\ \vdots \\ \boldsymbol{s}_n^T \end{bmatrix}}_{\text{3D structure matrix}} \underbrace{\begin{bmatrix} \boldsymbol{w}_1 & \boldsymbol{w}_2 & \cdots & \boldsymbol{w}_{2m} \end{bmatrix}}_{\text{camera motion matrix}}$$

where  $s_i \in \mathbb{R}^3$ , which gives rank(M) = 3.

Due to occlusions, there are many missing entries in the matrix M. Can we complete the missing entries?

Many more applications:

- spatial-temporal data: low-rank due to correlations, e.g. MRI video, network traffic, etc..
- quantum space tomography
- linear system identification

Problem of interest: Can we recover the matrices of interest from "incomplete" observations, using efficient algorithms?

• the problem is ill-posed without additional constraints

- Let  $M \in \mathbb{R}^{n \times n}$  (square case for simplicity) be a matrix of rank  $r \ll n$ .
- The Singular Value Decomposition (SVD) of M is given as

$$oldsymbol{M} = \sum_{i=1}^r \sigma_i oldsymbol{u}_i oldsymbol{v}_i^T$$

where  $\{\sigma_i\}_{i=1}^r$  are the singular values; and  $\{u_i\}_{i=1}^r$ ,  $\{v_i\}_{i=1}^r$  are the singular vectors.

• M has (2n-r)r degrees of freedom.

We make linear measurements of M:

$$y_i = \langle \boldsymbol{A}_i, \boldsymbol{M} \rangle = \operatorname{Tr}(\boldsymbol{A}_i^T \boldsymbol{M}), \qquad i = 1, \dots m,$$

which can be written more concisely in an operator form:

$$oldsymbol{y} = \mathcal{A}(oldsymbol{M})$$

where  $\mathcal{A} : \mathbb{R}^{n \times n} \mapsto \mathbb{R}^m$  denotes the measurement process. Its adjoint operator  $\mathcal{A}^* : \mathbb{R}^m \mapsto \mathbb{R}^{n \times n}$  is defined as

$$\mathcal{A}^*(\boldsymbol{y}) = \sum_{i=1}^m y_i \boldsymbol{A}_i.$$

The problem of rank minimization:

$$\hat{M} = \operatorname*{argmin}_{X} \operatorname{rank}(X) \quad \text{s.t.} \quad y = \mathcal{A}(X).$$

Just as  $\ell_1$  norm provides a convex relaxation to cardinality minimization, we use the nuclear norm which give a convex relaxation to rank minimization.

**Definition 1.** The nuclear norm of X is defined as

$$\|\boldsymbol{X}\|_* = \sum_{i=1}^n \sigma_i(\boldsymbol{X})$$

where  $\sigma_i(\mathbf{X})$  is the *i*th largest singular value of  $\mathbf{X}$ .

- Since the rank is  $\sum_{i=1}^{n} 1(\sigma_i(X) \neq 0)$ , the nuclear norm can be thought as an  $\ell_1$  norm relaxation of the vector of singular values.
- This is a norm. Relationships between different norms:

$$\|\boldsymbol{X}\| \leq \|\boldsymbol{X}\|_F \leq \|\boldsymbol{X}\|_* \leq \sqrt{r} \|\boldsymbol{X}\|_F \leq r \|\boldsymbol{X}\|.$$

• Tightest convex relaxation:  $\{X : \|X\|_* \le 1\}$  is the convex hull of rank-1 matrices obeying  $\|xy^T\| \le 1$ .

**Lemma 1.** Let A and B be matrices of the same dimensions. If  $AB^T = 0$ and  $A^TB = 0$ , then  $||A + B||_* = ||A||_* + ||B||_*$ .

Remark: this implies that, if the row and column spaces of A and B are orthogonal, then  $||A + B||_* = ||A||_* + ||B||_*$ .

This is similar to the  $\ell_1$  norm when  $m{x}$  and  $m{y}$  have disjoint support:

$$\|m{x} + m{y}\|_1 = \|m{x}\|_1 + \|m{y}\|_1$$

which is essentially all we need to get the proof of  $\ell_1$  minimization with RIP...

#### Lemma 2.

$$\|\boldsymbol{X}\|_{*} = \min_{\boldsymbol{W}_{1}, \boldsymbol{W}_{2}} \left\{ \frac{1}{2} \operatorname{Tr}(\boldsymbol{W}_{1}) + \frac{1}{2} \operatorname{Tr}(\boldsymbol{W}_{2}) \middle| \begin{bmatrix} \boldsymbol{W}_{1} & \boldsymbol{X} \\ \boldsymbol{X}^{T} & \boldsymbol{W}_{2} \end{bmatrix} \succeq 0 \right\}.$$

This means we can compute the nuclear norm efficient via semidefinite programming (SDP).

**Proof:** on the blackboard.

**Definition 2.** For a given norm  $\|\cdot\|_{\mathcal{A}}$  in an inner product space  $\langle\cdot,\cdot\rangle$ , the dual norm is defined as

$$\|\boldsymbol{X}\|_{\mathcal{A}}^{\star} := \max\{\langle \boldsymbol{X}, \boldsymbol{Y} \rangle : \|\boldsymbol{Y}\|_{\mathcal{A}} \leq 1\}.$$

By definition, this gives a general version of Cauchy-Schwarz inequality:

 $\langle \boldsymbol{X}, \boldsymbol{Y} 
angle \leq \| \boldsymbol{X} \|_{\mathcal{A}} \| \boldsymbol{Y} \|_{\mathcal{A}}^{\star}.$ 

Examples:

- The dual norm of  $\|\cdot\|_F$  is  $\|\cdot\|_F$ ;
- The dual norm of  $\|\cdot\|_1$  is  $\|\cdot\|_\infty$ ;
- The dual norm of  $\|\cdot\|_*$  is  $\|\cdot\|_;$

rank minimization vs cardinality minimization:

parsimony concept	cardinality	rank	
Hilbert Space norm	Euclidean	Frobenius	
sparsity inducing norm	$\ell_1$	nuclear	
dual norm	$\ell_{\infty}$	operator	
norm additivity	disjoint support	orthogonal row and column spaces	
convex optimization	linear programming	semidefinite programming	

Table 1: A dictionary relating the concepts of cardinality and rank minimization.

• The rank minimization problem:

$$\hat{oldsymbol{M}} = \operatorname*{argmin}_{oldsymbol{X}} \mathsf{rank}(oldsymbol{X}) \quad \mathsf{s.t.} \quad oldsymbol{y} = \mathcal{A}(oldsymbol{X}).$$

• We pose the following nuclear norm minimization algorithm:

$$\hat{M} = \operatorname*{argmin}_{X} \|X\|_{*}$$
 s.t.  $y = \mathcal{A}(X)$ ,

which can be solved efficiently via SDP:

$$\hat{\boldsymbol{M}} = \operatorname*{argmin}_{\boldsymbol{X}, \boldsymbol{W}_{1}, \boldsymbol{W}_{2}} \frac{1}{2} \operatorname{Tr}(\boldsymbol{W}_{1}) + \frac{1}{2} \operatorname{Tr}(\boldsymbol{W}_{2})$$
  
s.t.  $\boldsymbol{y} = \mathcal{A}(\boldsymbol{X}), \quad \begin{bmatrix} \boldsymbol{W}_{1} & \boldsymbol{X} \\ \boldsymbol{X}^{T} & \boldsymbol{W}_{2} \end{bmatrix} \succeq 0.$ 

• If  $\mathcal{A}$  satisfies the restricted isometry property for low-rank matrices:

**Definition 3.** The operator A satisfies the RIP of rank-r, if for any rank-r matrix, we have

$$(1-\delta_r) \|\boldsymbol{X}\|_F^2 \le \|\boldsymbol{\mathcal{A}}(\boldsymbol{X})\|_F^2 \le (1+\delta_r) \|\boldsymbol{X}\|_F^2$$

for  $0 \leq \delta_r \leq 1$ .

- If  $\{A_i\}_{i=1}^m$  are composed of i.i.d. Gaussian entries, then it satisfies the matrix RIP of order r with high probability, as soon as  $m \gtrsim nr$ .
- This allows us to develop almost parallel results to compressed sensing.

**Theorem 3.** If A satisfies the RIP of rank 4r with  $\delta_{4r} \leq \sqrt{2} - 1$ , then for all rank-r matrices, the nuclear norm minimization algorithm recovers M exactly.

Exact recovery from O(nr) measurements!!

• For the noisy case,

$$\boldsymbol{y} = \mathcal{A}(\boldsymbol{M}) + \boldsymbol{w}$$

where w is composed of i.i.d.  $\mathcal{N}(0, \sigma^2)$  entries. We could similarly propose the matrix LASSO algorithm:

$$\hat{\boldsymbol{M}} = \operatorname*{argmin}_{\boldsymbol{X}} \frac{1}{2} \|\boldsymbol{y} - \mathcal{A}(\boldsymbol{X})\|_{2}^{2} + \lambda \|\boldsymbol{X}\|_{*},$$

where  $\lambda$  is a regularization parameter.

• If  $\|\mathcal{A}^*(\boldsymbol{w})\| \leq \lambda/2$  and  $\delta_{4r} < (3\sqrt{2}-1)/17$ , then

$$\|\hat{\boldsymbol{M}} - \boldsymbol{M}\|_F \le C\sqrt{r}\lambda$$

for some constant C. For the Gaussian case,

$$\|\mathcal{A}^*(\boldsymbol{w})\| \le c_1 \sqrt{n} \sigma := \lambda$$

for some large enough constant  $c_1$  with probability at least  $1 - 2e^{-cn}$ .

• If M is an approximately low-rank matrix, we further have

$$\|\hat{\boldsymbol{M}} - \boldsymbol{M}\|_F \le C_1 \frac{\|\boldsymbol{M} - \boldsymbol{M}_r\|_*}{\sqrt{r}} + C_2 \sqrt{nr}\sigma$$

with probability at least  $1 - 2e^{-cn}$ , in the Gaussian sampling case.

In the matrix completion setting, we are given partial observations of the entries of M, and wish to recover the missing entries.

- Denote  $\Omega = \{(i, j) \in [n] \times [n]\}$  as the index set of observed entries.
- The observation can be written as

$$oldsymbol{Y} = \mathcal{P}_{\Omega}(oldsymbol{M})$$

where  $Y_{ij} = M_{ij}$  if  $(i, j) \in \Omega$  and  $Y_{ij} = 0$  otherwise.

• Consider the following algorithm:

$$\min \|\boldsymbol{X}\|_*$$
 s.t.  $\boldsymbol{Y} = \mathcal{P}_{\Omega}(\boldsymbol{X})$ 

• The observation operator doesn't satisfy RIP!

Consider a rank-one matrix  $M = xy^T$  with the following sampling pattern:

×	×	×	×	×	×
×	×	×	×	×	×
×	×	$\times$	$\times$	$\times$	$\times$
×	×	$\times$	$\times$	$\times$	$\times$
L×	$\times$	×	×	×	×

If single row (or column) is not sampled, recovery is not possible.

Fix the number of observed entries  $m = |\Omega|$ , would like to get performance bound that holds for almost all sampling patterns.

 $\implies$  We'll consider subset of m entries selected uniformly at random.

Compare the following two rank-one matrices:

$$\begin{bmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & & \\ 0 & 0 & 0 & \cdots & 0 \end{bmatrix} \quad \mathbf{vs} \quad \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \vdots & & \\ 1 & 1 & 1 & \cdots & 1 \end{bmatrix} \quad \mathbf{vs} \quad \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 0 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & & \\ 0 & 0 & 0 & \cdots & 0 \end{bmatrix}$$

The middle one would be "easier" to complete.

Column and row spaces cannot be aligned with basis vectors.

Let M be a rank-r matrix with the SVD  $M = U\Sigma V^T$ , where  $U, V \in \mathbb{R}^{n \times r}$ .

**Definition 4.** [Coherence] Smallest scalar  $\mu$  obeying

$$\max_{1 \le i \le n} \| \boldsymbol{U}^T \boldsymbol{e}_i \|_2^2 \le \mu \frac{r}{n}, \quad \max_{1 \le i \le n} \| \boldsymbol{V}^T \boldsymbol{e}_i \|_2^2 \le \mu \frac{r}{n},$$

where  $e_i$  is the *i*th standard basis vector.



- Geometric condition:  $\boldsymbol{U} = \operatorname{colspan}(\boldsymbol{M})$
- Since  $\sum_{i=1}^{n} \| \boldsymbol{U}^{T} \boldsymbol{e}_{i} \|_{2}^{2} = r$ ,  $\mu \geq 1$ .
- If  $oldsymbol{e}_i\inoldsymbol{U}$ ,  $\mu=n/r$ ;

• If 
$$\frac{1}{\sqrt{n}}\mathbf{1} = \boldsymbol{U}$$
,  $\mu = 1$ .

We would like  $\mu = O(1)$ .

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### Information-theoretic lower bound

**Theorem 4.** [Candes and Tao, 2009] No method can succeed with

 $m \lesssim \mu \times nr \times \log n \approx \textit{dof} \times \mu \log n$ 

Remarks:

- When  $\mu = O(1)$ , we need  $m \leq nr \log n$ .
- Need at least one observation /row and column related to the coupon collector's problem: Suppose that there is an urn of n different coupons, from which coupons are being collected, equally likely, with replacement. How many trials do we need to collect all n coupons?
- The adjacency graph needs to be fully connected

**Theorem 5.** [Chen, Gross, Recht, Candes and Tao, etc..] There exists universal constant  $c_0, c_1, c_2 > 0$  such that if

 $m \ge c_0 \mu nr \log^2 n,$ 

then M is the unique optimal solution to the nuclear norm minimization problem with probability at least  $1 - c_1 n^{-c_2}$ .

#### Remark:

• This result is optimal up to a logarithmic factor in *n*. See [Chen, Incoherence-Optimal Matrix Completion].

### Geometry



- $oldsymbol{P}_U$  is the orthogonal projection to the column space of  $oldsymbol{M}$ ;
- $oldsymbol{P}_V$  is the orthogonal projection to the row space of  $oldsymbol{M}$ ;
- Let T be the span of matrices of the form:

$$T = \{ \boldsymbol{U}\boldsymbol{X}^T + \boldsymbol{Y}\boldsymbol{V}^T : \boldsymbol{X}, \boldsymbol{Y} \in \mathbb{R}^{n \times r} \}$$

• Let  $\mathcal{P}_T$  be the orthogonal projection onto T:

$$\mathcal{P}_T(\boldsymbol{X}) = \boldsymbol{P}_U \boldsymbol{X} + \boldsymbol{X} \boldsymbol{P}_V - \boldsymbol{P}_U \boldsymbol{X} \boldsymbol{P}_V$$

• The complement projection  $\mathcal{P}_{T^{\perp}} = \mathcal{I} - \mathcal{P}_{T}$ :

$$\mathcal{P}_{T^{\perp}}(\boldsymbol{X}) = (\boldsymbol{I} - \boldsymbol{P}_U)\boldsymbol{X}(\boldsymbol{I} - \boldsymbol{P}_V)$$

The subgradient of  $\|\cdot\|_*$  at M can be written as

$$\partial \| \boldsymbol{M} \|_* = \left\{ \boldsymbol{U} \boldsymbol{V}^T + \boldsymbol{W} : \quad \mathcal{P}_T(\boldsymbol{W}) = 0, \| \boldsymbol{W} \| \le 1 \right\}$$

 $oldsymbol{Z}\in\partial\|oldsymbol{M}\|_*$  if and only if

$$\mathcal{P}_T(\boldsymbol{Z}) = \boldsymbol{U}\boldsymbol{V}^T, \quad \|\mathcal{P}_{T^{\perp}}(\boldsymbol{Z})\| \leq 1.$$

The subgradient doesn't depend on the singular values of M.

For any  $(i, j) \in [n] \times [n]$ :

$$\left\| \mathcal{P}_T(\boldsymbol{e}_i \boldsymbol{e}_j^T) \right\|_F^2 \leq \frac{2\mu r}{n}.$$

The sampling basis is incoherent to the tangent space T.

It turns out it is easier to use a sampling with replacement model, where we assume each observed entry  $(i_k, j_k)$ ,  $k = 1, \ldots, m$  is **i.i.d.** observed uniformly at random from  $[n] \times [n]$ .

This is much easier to analyze, however it is different from the sampling without replacement model stated earlier because we may sample the same entry several times.

**Proposition 1.** The probability that the nuclear norm heuristic fails when the set of observed entries is sampled uniformly from the collection of sets of size m is less than or equal to the probability that the heuristic fails when m entries are sampled independently with replacement.

**Proposition 2.** With probability at least  $1 - n^{2-2\beta}$ , the maximum number of repetitions of any entry in  $\Omega$  is less than  $\frac{8}{3}\beta \log(n)$  for  $n \ge 9$  and  $\beta > 1$ .

Define the operator

$$\mathcal{R}_{\Omega}(\boldsymbol{X}) = \sum_{k=1}^{m} \left\langle \boldsymbol{X}, \boldsymbol{e}_{i_k} \boldsymbol{e}_{j_k}^T \right\rangle \boldsymbol{e}_{i_k} \boldsymbol{e}_{j_k}^T = \sum_{k=1}^{m} X_{i_k, j_k} \boldsymbol{e}_{i_k} \boldsymbol{e}_{j_k}^T$$

where  $(i_k, j_k)$  is uniformly drawn from  $[n] \times [n]$ . From the above proposition, we have

$$\|\mathcal{R}_{\Omega}\| \le \frac{8}{3}\beta\log(n)$$

with probability at least  $1 - n^{2-2\beta}$ .

**Proposition 3.** [Exact Dual Certificate] M is the unique minimizer of the nuclear norm minimization problem if the following holds:

- the sampling operator  $\mathcal{P}_{\Omega}$  restricted to elements in T is injective;
- there exists Y supported on  $\Omega$  such that  $Y \in \partial \|M\|_*$ , i.e.

$$\mathcal{P}_T(\boldsymbol{Y}) = \boldsymbol{U}\boldsymbol{V}^T, \quad \|\mathcal{P}_{T^{\perp}}(\boldsymbol{Y})\| \leq 1.$$

The first equality constraint is not easy to satisfy, see [Candes and Tao, 2009].

Under a stronger injectivity requirement, we can relax the second requirement a bit, which much simplifies the analysis.

**Proposition 4.** [Inexact Dual Certificate] Suppose that

$$\frac{n^2}{m} \left\| \mathcal{P}_T \mathcal{R}_\Omega \mathcal{P}_T - \frac{m}{n^2} \mathcal{P}_T \right\| \le \frac{1}{2},$$

and there exists  $\boldsymbol{Y}$  supported on  $\Omega$  such that

$$\left\| \mathcal{P}_T(\mathbf{Y}) - \mathbf{U}\mathbf{V}^T \right\|_F \le \sqrt{\frac{r}{2n}}, \quad \left\| \mathcal{P}_{T^{\perp}}(\mathbf{Y}) \right\| < \frac{1}{2},$$

then M is the unique minimizer of the nuclear norm minimization problem if the following holds:

**Proposition 5.** For all  $\beta > 1$ ,

$$\frac{n^2}{m} \left\| \mathcal{P}_T \mathcal{R}_\Omega \mathcal{P}_T - \frac{m}{n^2} \mathcal{P}_T \right\| \le \sqrt{\frac{32\beta\mu_0 nr\log n}{3m}}$$

with probability at least  $1 - 2n^{2-2\beta}$  provided  $m \ge \frac{32}{3}\beta\mu_0 nr\log n$ .

**Remark:** Provided 
$$\sqrt{\frac{32\beta\mu_0 nr\log n}{3m}} \leq \frac{1}{2}$$
, i.e.

$$m \ge \frac{128\beta\mu_0 nr\log n}{3}$$

we have with probability at least  $1-2n^{2-2\beta}$  ,

$$\frac{n^2}{m} \left\| \mathcal{P}_T \mathcal{R}_\Omega \mathcal{P}_T - \frac{m}{n^2} \mathcal{P}_T \right\| \le \frac{1}{2}.$$

We introduce the clever golfing scheme proposed by David Gross.



Fig. 3. Caricature of the "golfing scheme" used to construct the certificate. In the *i*th step,  $X_{i-1}$  designates the vector we aim to represent. The approximation of  $X_{i-1}$  actually obtained is  $\mathcal{P}_T \mathcal{R}_i X_{i-1}$ . The distance of the new goal  $X_i = X_{i-1} - \mathcal{P}_T \mathcal{R}_i X_{i-1}$  to the origin is guaranteed to be only half the previous one. The sequence  $X_i$  thus converges exponentially fast to the origin.

Recall the FISTA algorithm we discussed to solve

$$\hat{\boldsymbol{M}} = \operatorname*{argmin}_{\boldsymbol{X}} \frac{1}{2} \|\boldsymbol{y} - \boldsymbol{\mathcal{A}}(\boldsymbol{X})\|_{2}^{2} + \lambda \|\boldsymbol{X}\|_{*}$$

• Initialization: 
$$oldsymbol{x}_0 = oldsymbol{x}_{-1} \in \mathbb{R}^n$$
,  $oldsymbol{ heta}_0 = 1$ ,

• For k = 1, 2, ...,

$$\begin{aligned} \theta_k &= \frac{1 + \sqrt{1 + 4\theta_{k-1}^2}}{2} \\ \mathbf{Y}_k &= \mathbf{X}_{k-1} + \left(\frac{\theta_{k-1} - 1}{\theta_k}\right) \left(\mathbf{X}_{k-1} - \mathbf{X}_{k-2}\right) \\ \mathbf{X}_k &= \operatorname{prox}_{t_k \lambda \|\mathbf{X}\|_*} \left(\mathbf{Y}_k - t_k \mathcal{A}^* (\mathcal{A}(\mathbf{Y}_k) - \mathbf{y})\right) \end{aligned}$$

• What is the proximal operator for  $\|\cdot\|_*$ ?

**Proposition 6.** 

$$\operatorname{prox}_{t_k\lambda\|\cdot\|_*}(\boldsymbol{X}) = \operatorname{argmin}_{\boldsymbol{Z}} \left\{ \frac{1}{2} \|\boldsymbol{Z} - \boldsymbol{X}\|_2^2 + t_k\lambda\|\boldsymbol{Z}\|_* \right\} = \mathcal{T}_{t_k\lambda}(\boldsymbol{X})$$

where

$$\mathcal{T}_{\tau}(\boldsymbol{X}) = \boldsymbol{U}\mathcal{T}_{\tau}(\boldsymbol{\Sigma})\boldsymbol{V}^{T},$$

where the SVD of X is given as  $X = U\Sigma V^T$ ,  $\Sigma = diag(\{\sigma_i\})$ , and

$$\mathcal{T}_{\tau}(\mathbf{\Sigma}) = diag(\{(\sigma_i - \tau)_+\}).$$