

# Algorithms and Computation in Signal Processing

special topic course 18-799B  
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# Miscellaneous

- Assignments on Website, due next Tuesday before class
- PC for programming assignments
- Project

# Asymptotic Analysis (cont'd)

- Problems with fixed upper bound on input size
  - The complexity of multiplying 2 matrices of dimensions smaller than  $1000 \times 1000$  has complexity  $O(1)$
- Multiple parameters (blackboard)
  - Mat-mat multiplication
  - Polynomial multiplication

# Asymptotic Analysis: Remarks

- Asymptotic runtime analysis works since it is **independent** of the exact runtime of the elementary steps counted (including memory latencies)
- Complexity of a problem is usually stated using “O” since every algorithms provides an upper bound, but lower bounds are hard to get
- People often talk about “complexity of an algorithm” which, in a strict sense, is wrong
- Problem: asymptotic analysis gives only an asymptotic idea of the runtime, but in real implementations the constants (and the algorithmic structure) matter

# Cost Analysis

# Refined Analysis for Numerical Problems

- Goal: determine exact static “cost” of algorithms
- Approach (use mat-mat mult. as running example):
  - Fix an appropriate cost measure  $C$ , “what do I count”
  - Determine cost of algorithm as function  $C(n)$  of input size  $n$ , or, more general, of all relevant input parameters:
$$C(n_1, \dots, n_k)$$
  - Cost can be multi-dimensional
$$C(n_1, \dots, n_k) = (c_1, \dots, c_m)$$
- Exact cost gives a more precise idea of runtime but not the exact runtime

# Cost Analysis

## ■ Examples

- Count additions and multiplications for flop rate
- Mat-mat mult.
- Polynomial mult.

## ■ Solving recurrences

- Great book: Graham, Knuth, Patashnik, "Concrete Mathematics," 2<sup>nd</sup> edition, Addison Wesley 1994
- Blackboard

# For Publications

- A problem has a complexity
- An algorithm has a cost (e.g., operations count, runtime, memory requirement, area requirement in hardware)
- Cost=runtime can only be analyzed asymptotically
- In a precise sense, an algorithm does not have a complexity

Problem	Complexity Runtime compl. (asympt.)
Algorithm	Cost Runtime (asymptotic)

## In research/writing/publications:

If your contribution is an algorithm, you have to analyze it. As follows:

- 1) state your cost/complexity measure (what you count);
- 2) compute the cost of the algorithm as precise as possible/necessary, at least asymptotically;
- 3) state what you know about the complexity of the problem you address (from theory, other algorithms, ...)