

Building Safer UGVs with Run-time Safety Invariants

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October 28, 2009

UNCLASSIFIED: Distribution A. Approved for Public Release TACOM Case #20247 Release Date: 07 OCT 2009

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The Message



- To be useful, unmanned ground vehicles (UGVs) must safely operate alongside personnel, although this is not yet reliable enough with today's technology.
- The use of physical safety barriers and large stand-off distances is acceptable only during testing; it is infeasible for use in the real world.
- We are developing safeguards to reduce dependence on physical barriers and large standoff distances for UGV operating alongside personnel in real, dynamic operations.



- Our approach is based on *run-time* safety *invariants* enforced by a Safety Monitor
- Benefits of our approach involve
 - →A clear definition of "safety"
 - →Firewalling safety-criticality to a small set of components
 - →Streamlined V&V of safety-critical components
- We are implementing our approach on the Autonomous Platform Demonstrator project
- We will discuss our process for developing a Safety Monitor using the Autonomous Platform Demonstrator (APD) as an example



- Run-time safety invariants are concise, formal expressions of critical system properties that define system safety
 - \rightarrow E.g., "vehicle speed doesn't exceed operator-specified limit"
 - \rightarrow We needn't enumerate detailed causes of hazards
 - →Rather, we create a dependable outer bound on what it means to be "safe"
 - \rightarrow Do this based on fault-tree analysis

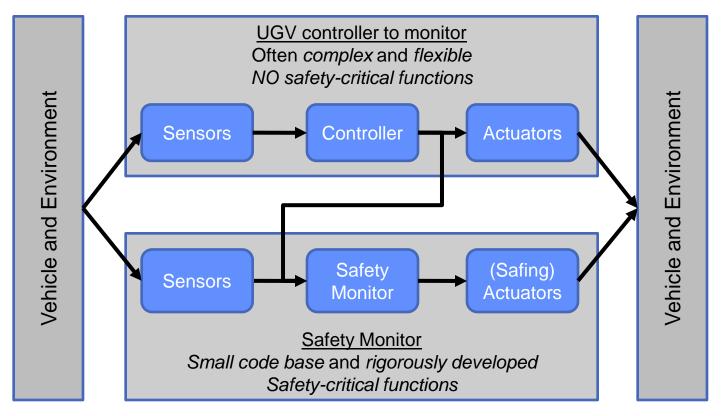
Our Approach (2)



• We then build a Safety Monitor that safes the UGV whenever any invariant is violated

 \rightarrow Has a dependable means of sensing invariant state

 \rightarrow Has a dependable means of safing the system



Demonstration Vehicle: APD



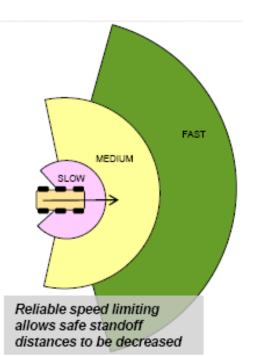


APD Safety Goals



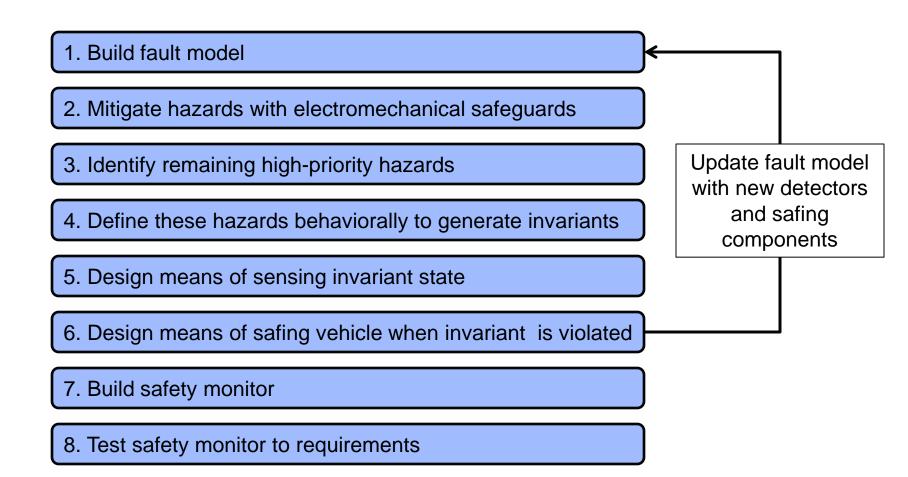
• Initial focus is on mitigating hazards involved with driving the APD vehicle

- →Ensure the vehicle can be stopped when commanded
- →Ensure the vehicle maintains a commanded speed limit
- Meeting both these goals helps to decrease safe standoff distances



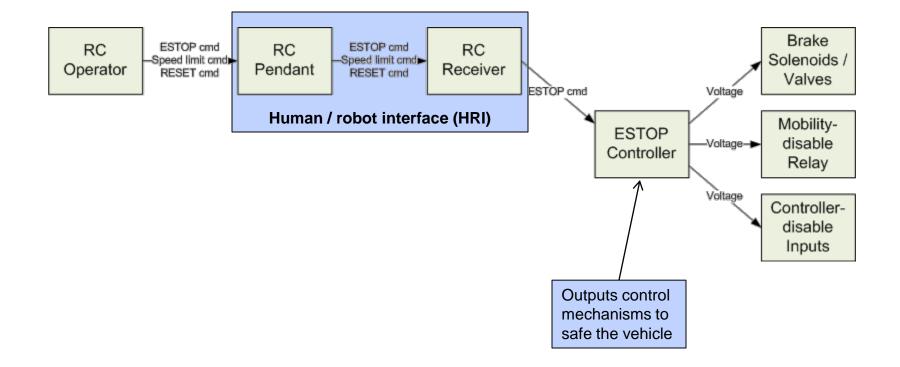
Development Process





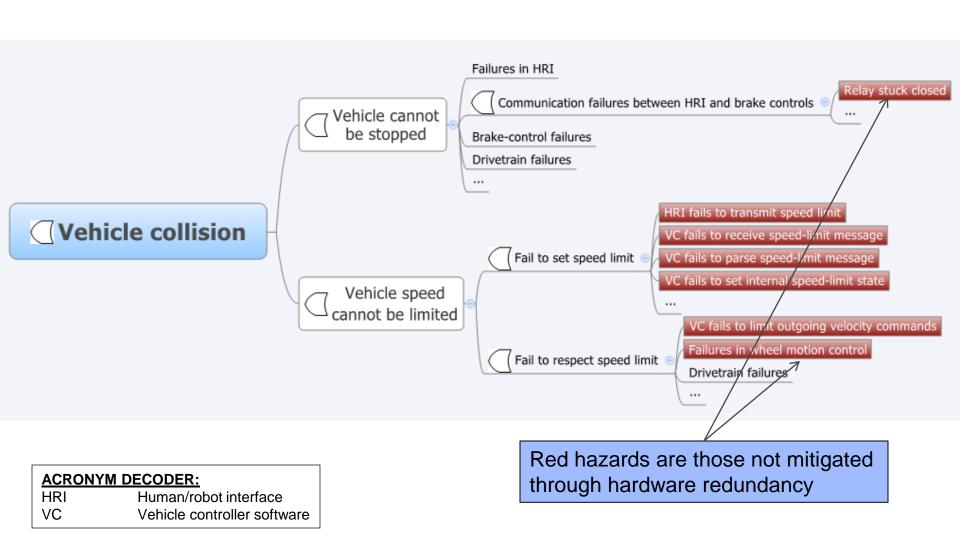
APD Safety Architecture





ACRONYM DECODER:	
ESTOP	Emergency stop
RC	Radio controller

APD Fault Model Example





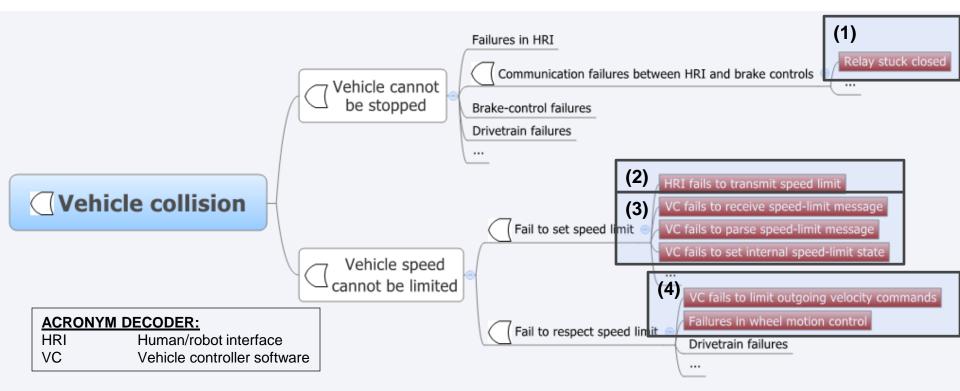
Hazard	Behavioral Definition
(HRI) relay stuck closed	HRI reports ESTOP signal over serial line but relay is closed
Failures in HRI	No valid heartbeat from HRI
Communication failures between HRI and VC	No valid heartbeat from HRI No valid heartbeat from VC
VC fails to parse speed-limit message	Vehicle exceeds speed limit specified by HRI
VC fails to set internal speed-limit state	
VC fails to limit outgoing velocity commands	
Failures in wheel motion control	

APD Safety Invariants



Safe the vehicle if:

- 1. HRI ESTOP is commanded, OR
- 2. HRI is inactive, OR
- 3. VC is inactive, OR
- 4. Vehicle speed exceeds limit specified by HRI



Means of Sensing Invariants



1. HRI ESTOP command

- \rightarrow Data packets received from HRI
- \rightarrow Packets include error-detection code
- 2. HRI is inactive
 - \rightarrow Valid packet received from HRI
- 3. VC is inactive
 - → Valid driving command received sent by VC and snooped by Safety Monitor

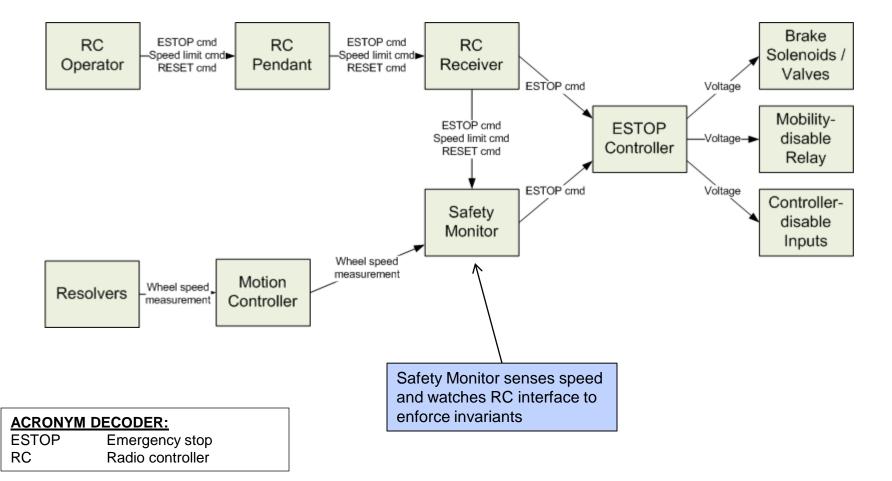
4. <u>Vehicle speeds exceed limit specified by HRI</u>

- → Wheel velocities are reported through telemetry from lowlevel traction drive controllers
- \rightarrow Data packets from HRI specify setting of a speed-limit switch



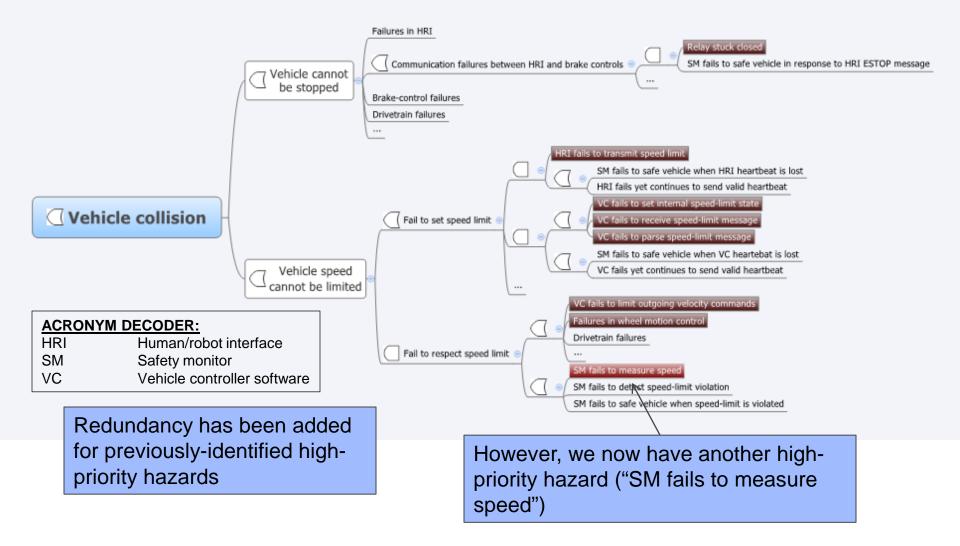
- Must be...
 - \rightarrow Independent of non-safety critical components
 - \rightarrow Unable to be overridden or disabled
 - →Fail-safe
- On APD, an ESTOP-controller applies fail-safe mechanical brakes if any of a set of inputs drop low
- The safety monitor has control over one of these inputs

Updated Safety Architecture



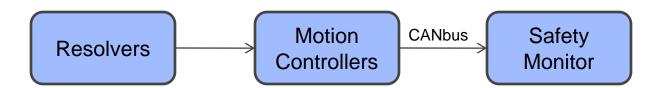
Updated Fault Model





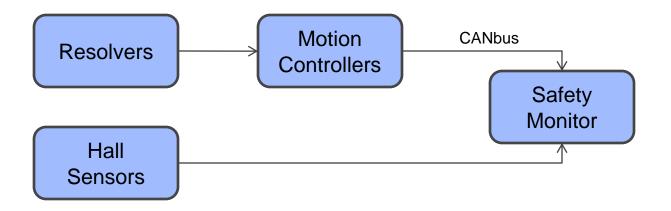


- Speed is sensed through telemetry from motioncontrol hardware
 - →Vehicle speed is estimated as an average of wheel speeds
- These motion controllers are "black boxes" supplied by a vendor, so thorough V&V is infeasible
 - →Control hardware could report false readings
 - → Firmware changes could have unintended consequences
 - →Resolvers could fail





- To address these risks we added redundant wheelspeed sensing
- Hall-effect sensors are placed in hubs that are wired directly to the safety monitor
- Use these sensors to check the validity of measurements from the motion controllers



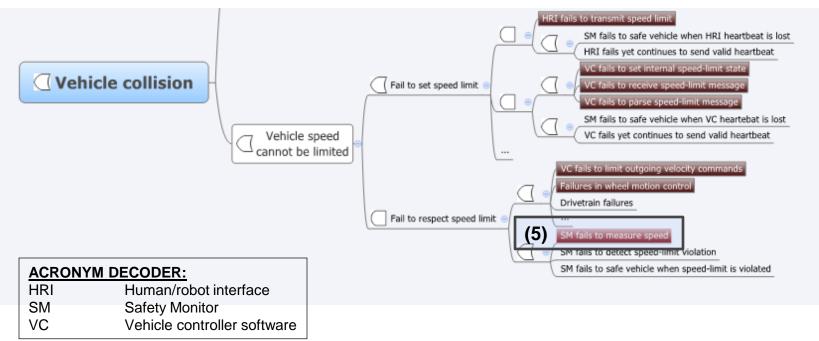


- A failure of one sensing modality will not affect readings from the other:
 - →Largely separate power supplies
 - →Motion control firmware completely separate from hall sensors
 - →Motion controllers communicate via CAN bus, hall sensors use separate dedicated inputs
 - $\rightarrow \mbox{Resolvers}$ and hall sensors mounted in different locations



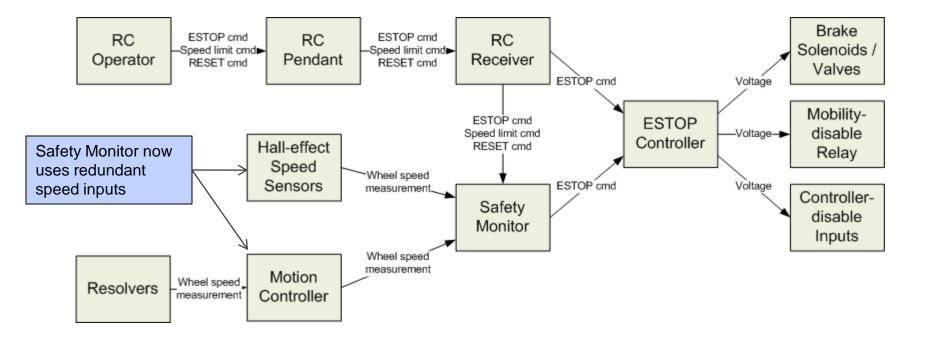
Safe the vehicle if:

- 1. HRI ESTOP is commanded, OR
 - 2. HRI is inactive, OR
 - 3. VC is inactive, OR
 - 4. Vehicle speed exceeds limit specified by HRI, OR
 - 5. Vehicle-speed measurements disagree



Final Architecture

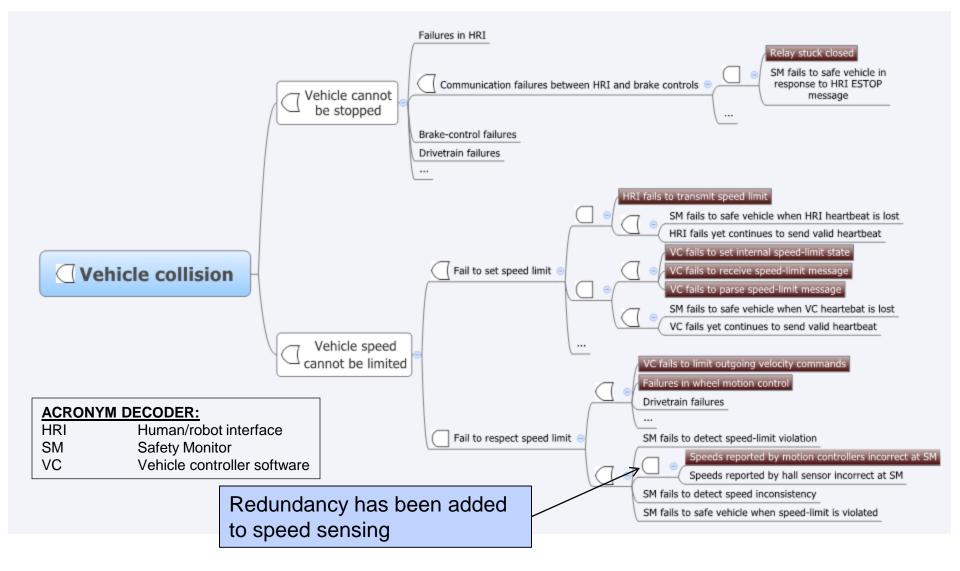




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Final Fault Model

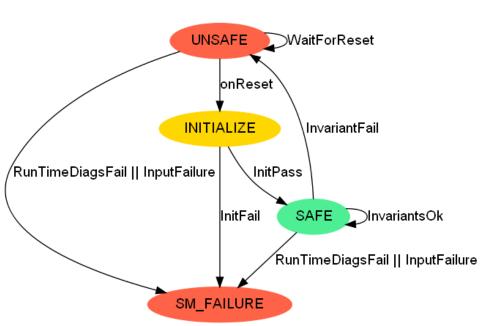




Safety Monitor Design



- Simple finite state machine design
- If an invariant is violated, enter UNSAFE state and trigger ESTOP
 - \rightarrow Return to SAFE state once invariants again hold and operator issues RESET
- If any self-checks fail, assume SM cannot evaluate invariants
 - →Enter SM_ASSERT state, which halts execution and triggers ESTOP with an independent hardware watchdog



Safety Monitor Master State Chart

Safety Monitor Implementation



 Implement as a single work loop

→Minimize use of interrupt I/O as much as possible

- Separate processing of input sources (e.g., conversion of hall-sensor readings to vehicle speed) from invariant evaluation
- Evaluate invariants based on simple boolean functions

while (true)

process_input_data()
evaluate_invariants()
update_SM_state()
set_ESTOP_output()
send_status_output()



- The approach results in simpler test goals than we'd have if we had to verify a complex safety system
 - \rightarrow 80% of project resources are typically spent on V&V
 - \rightarrow So streamlining V&V results in bigger payoffs than improving development tools
- Safety invariants are testable safety requirements
- For each invariant, carry out:
 - \rightarrow <u>System test</u> that the SM issues an ESTOP if the invariant is violated
 - $\rightarrow \underline{\text{Bench test}}$ that the SM issues an ESTOP if invalid input signals are received
 - \rightarrow <u>Unit test</u> that the SM transitions to UNSAFE state upon any time-based combination of invariant-violation
 - → <u>Code review</u> that the processing of input data for the evaluation of invariants is correct
- Prove and Document that the means of safing the system is fail-safe

References



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